

Part 1 triangulated categories

roughly! exact category / homotopy
→ want notion of long exact sequence

Definition A triangulated category consists of

- an additive category \mathcal{C}
- an autoequivalence $\Sigma: \mathcal{C} \rightarrow \mathcal{C}$ called suspension
- a class of ~~exact~~ triangles Δ called exact triangles

A triangle is a diagram of the form

$$A \xrightarrow{f} B \xrightarrow{g} C \xrightarrow{h} \Sigma A$$

They must satisfy

(TR1) • Δ is closed under isomorphisms
• Every morphism $A \xrightarrow{f} B$ in \mathcal{C} can be completed to an ~~exact~~ exact triangle $A \xrightarrow{f} B \rightarrow C \rightarrow \Sigma A$

(TR2) (rotation)

If $A \xrightarrow{f} B \xrightarrow{g} C \xrightarrow{h} \Sigma A$ is exact

then $B \xrightarrow{g} C \xrightarrow{h} \Sigma A \xrightarrow{\Sigma f} \Sigma B$ is exact

Note: iterated rotation yields a diagram of the form

$$\begin{array}{ccccccc} A & \rightarrow & B & \rightarrow & C & \rightarrow & \Sigma A \\ \hookrightarrow & & \hookrightarrow & & \hookrightarrow & & \hookrightarrow \\ \Sigma A & \rightarrow & \Sigma B & \rightarrow & \Sigma C & \rightarrow & \Sigma^2 A \\ \hookrightarrow & & \hookrightarrow & & \hookrightarrow & & \hookrightarrow \\ \Sigma^2 A & & & & & & \dots \end{array}$$

This allows construction of long exact sequences from triangles

(TR3) Given a diagram with exact rows

$$\begin{array}{ccccccc} A & \rightarrow & B & \rightarrow & C & \rightarrow & \Sigma A \\ \downarrow f & & \downarrow g & & \downarrow h & & \downarrow \Sigma f \\ A' & \rightarrow & B' & \rightarrow & C' & \rightarrow & \Sigma A' \end{array}$$

there is a morphism ψ making everything commute.

(TR4) The octahedral axiom
roughly: $A/C/B/C \cong A/C$

Example

Derived categories \rightarrow other talks

Stable category of Frobenius category \rightarrow this talk

Part 2 Frobenius Categories

Running example: Let A be l.d. algebra and
consider exact structures on $\text{mod-}A$:

S_{split} : conflations are split exact sequences

S_{short} : conflations are short exact sequences

Definition Let (\mathcal{F}, S) be an exact category.

$I \in \mathcal{F}$ is injective if $\mathcal{F}(-, I)$ maps conflations
to \rightarrow s.e.s.

\mathcal{F} has enough ^{injectives} projectives if every $A \in \mathcal{F}$ fits into
a conflation $A \rightarrow I \rightarrow \Sigma A$ with I injective

Duality: Projective, enough projectives

Example

For $(\text{mod-}A, S_{\text{split}})$ all objects are projective and injective

For $(\text{mod-}A, S_{\text{short}})$ projectives and injectives are the
usual ones.

Definition (\mathcal{F}, S) is a Frobenius category if

- \mathcal{F} has enough projectives and injectives
- The projectives and injectives coincide

Example $(\text{mod-}A, \mathcal{S}_{\text{split}})$ is Frobenius

$(\text{mod-}A, \mathcal{S}_{\text{short}})$ is Frobenius iff A is self-injective

Definition Let $(\mathcal{F}, \mathcal{S})$ be a Frobenius category.

The stable category $\underline{\mathcal{F}}$ has

• objects: same as \mathcal{F}

• morphisms: $\underline{\mathcal{F}}(A, B) := \mathcal{F}(A, B) / \left(\begin{array}{l} \text{morphisms that} \\ \text{factor through} \\ \text{projectives} \end{array} \right)$

\rightarrow A notion of homotopy

\rightarrow expect $\underline{\mathcal{F}}$ triangulated

Note: Let $P \in \mathcal{F}$ be projective then $Op = \text{id}_P$ in $\underline{\mathcal{F}}$

$\Rightarrow P \cong 0$ in $\underline{\mathcal{F}}$

$\rightarrow (\text{mod-}A, \mathcal{S}_{\text{split}})$ is trivial

$(\text{mod-}A, \mathcal{S}_{\text{short}})$ is more interesting

Theorem Let $(\mathcal{F}, \mathcal{S})$ be a Frobenius category

then $\underline{\mathcal{F}}$ is a triangulated category.

Construction: need \mathcal{S} and A

\mathcal{S} : for every $A \in \mathcal{F}$ Pick a conflation

$A \rightarrow I \rightarrow \mathcal{S}A$ with I injective.

Schanuel's Lemma $\Rightarrow \mathcal{S}A$ is unique up to ~~the~~ injective

direct summands $\rightarrow \mathcal{S}A$ is unique up to iso in $\underline{\mathcal{F}}$

\rightarrow can turn \mathcal{S} into a functor

pick a second conflation $\mathcal{S}A \rightarrow P \rightarrow A$ with P projective.

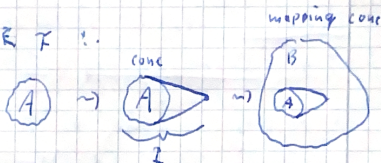
\rightarrow can turn \mathcal{S} into a quasi-inverse of \mathcal{S}

Note: The conflation $A \rightarrow I \rightarrow \Sigma A$ shows

$A \cong \Omega \Sigma A$ in \mathcal{E} since I is projective

Δ : Let $f: A \rightarrow B$ & $F: \dots$

topological motivation:



Construction: Let $A \xrightarrow{L} I \xrightarrow{P} \Sigma A$ be a conflation with I injective. Consider $B \oplus_A I$.

obtain exact triangle

$$A \xrightarrow{f} B \xrightarrow{(\partial, 0)} B \oplus_A I \xrightarrow{(0, P)} \Sigma A \quad \square$$

Part 3 Let \mathcal{A} be an abelian category.

Definition A full subcategory \mathcal{W} of \mathcal{A} is wide if

- it is closed under summands
- if two objects in an s.e.s. are in \mathcal{W} then so is the third

Definition Let \mathcal{X} be a full subcategory of \mathcal{A} then

$$\mathcal{X}^\perp := \{Y \in \mathcal{A} \mid \text{Ext}^1(X, Y) = 0\}$$

$${}^\perp \mathcal{X} := \{Y \in \mathcal{A} \mid \text{Ext}^1(Y, X) = 0\}$$

A pair $(\mathcal{X}, \mathcal{Y})$ of full subcategories is a cotorsion pair

if $\mathcal{X}^\perp = \mathcal{Y}$ and $\mathcal{X} = {}^\perp \mathcal{Y}$

it is functorially complete if every $A \in \mathcal{A}$ admits ^{functorial}

s.e.s. $0 \rightarrow Y \rightarrow X \rightarrow A \rightarrow 0$ and $0 \rightarrow A \rightarrow Y' \rightarrow X' \rightarrow 0$

with $X, X' \in \mathcal{X}$ and $Y, Y' \in \mathcal{Y}$

it is hereditary if $\text{Ext}^2(X, Y) = 0$